

Alternative Synchronization Strategies — Lock-Free Algorithms (1) —

lecture 19 (2021-05-24)

Master in Computer Science and Engineering

— Concurrency and Parallelism / 2020-21 —

João Lourenço <joao.lourenco@fct.unl.pt>

Alternative Synchronization Strategies

- Contents:

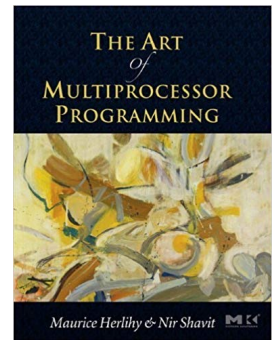
- Liveness: Types of Progress
- Coarse-Grained Synchronization
- Fine-Grained Synchronization
- Optimistic Synchronization
- Lazy Synchronization
- Lock-Free Synchronization

Past lectures

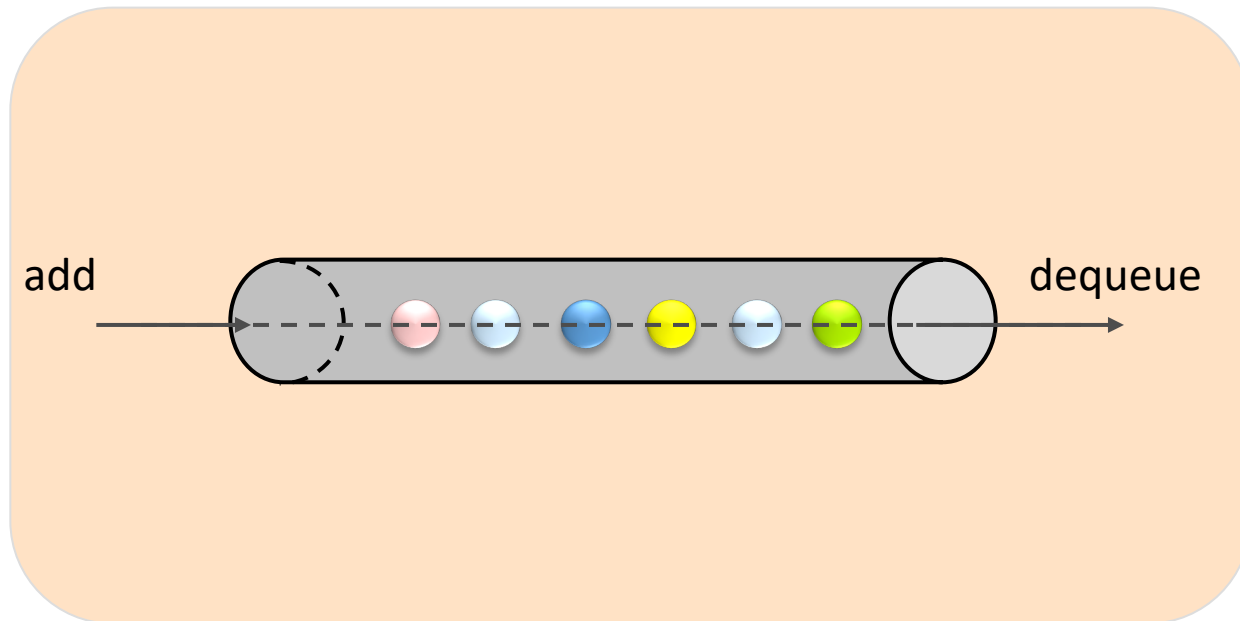
Today

- Reading list:

- chapter 5 of the Textbook
- Chapter 9 of “The Art of Multiprocessor Programming” by Maurice Herlihy & Nir Shavit *(available at clip)*

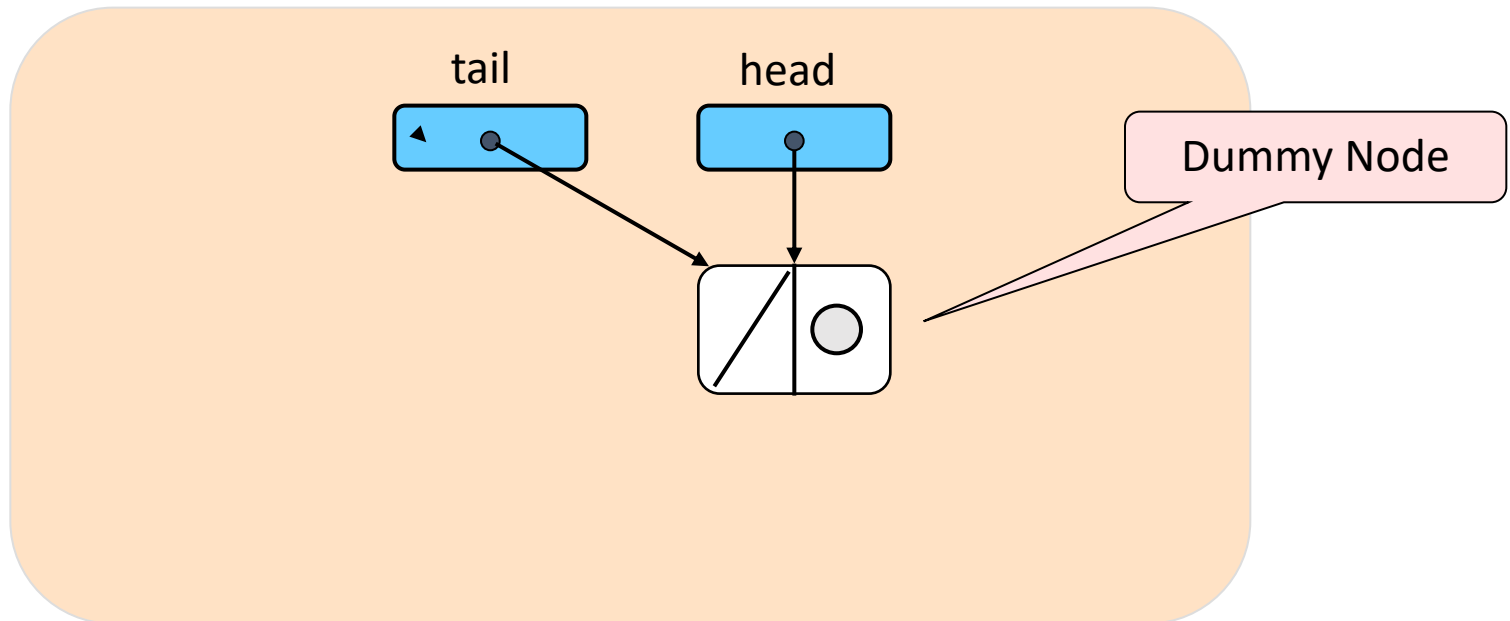


Basics for a lock-free Queue

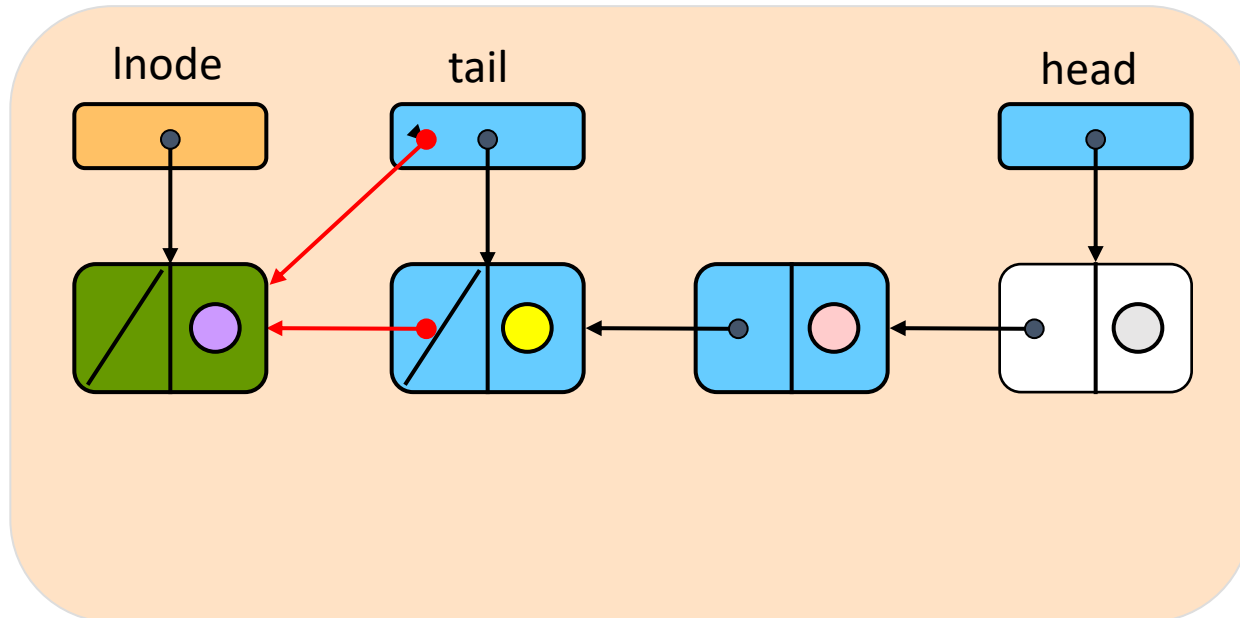


Basics for a lock-free Queue

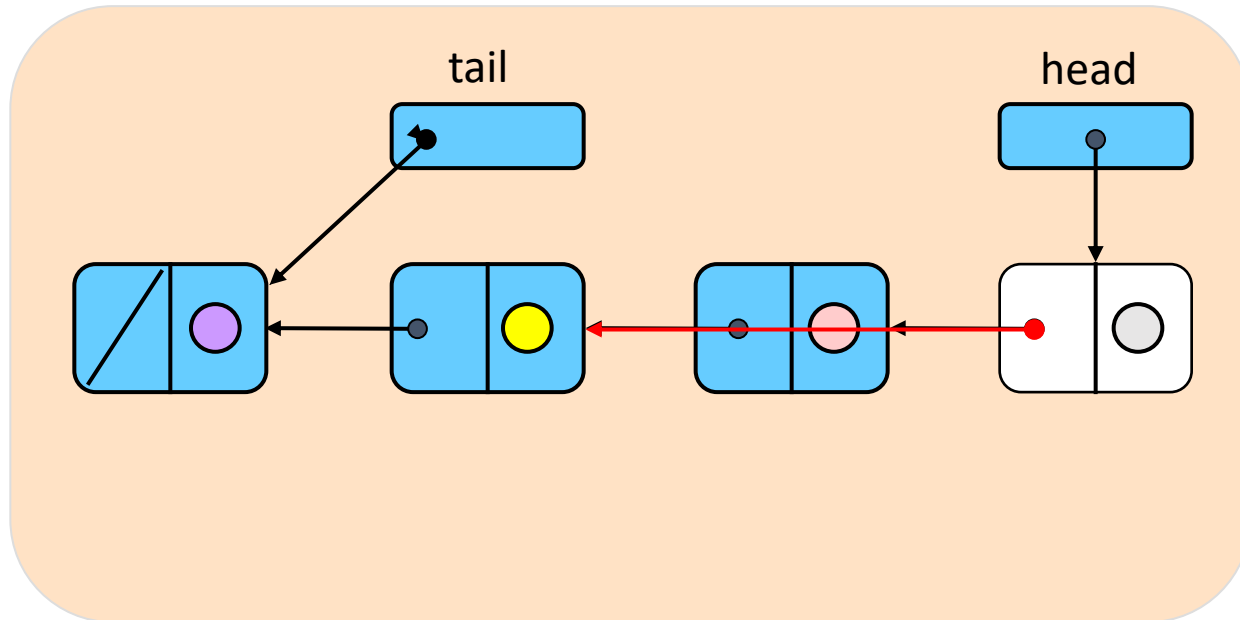
Empty queue



Enqueue



Deque



Compare & set(CAS)

shared
register old new


↓ ↓ ↓

CAS (A, B, C)

```
if A==B then A←C; return(true)
           else return(false)
```

Supported by Intel, AMD, Arm, ...

Reminder: Lock-Free Data Structures

- No matter what ... 
- Guarantees minimal progress in any execution
 - i.e., some thread will always complete a method call
- Even if others halt at malicious times
- Implies that implementation can't use locks

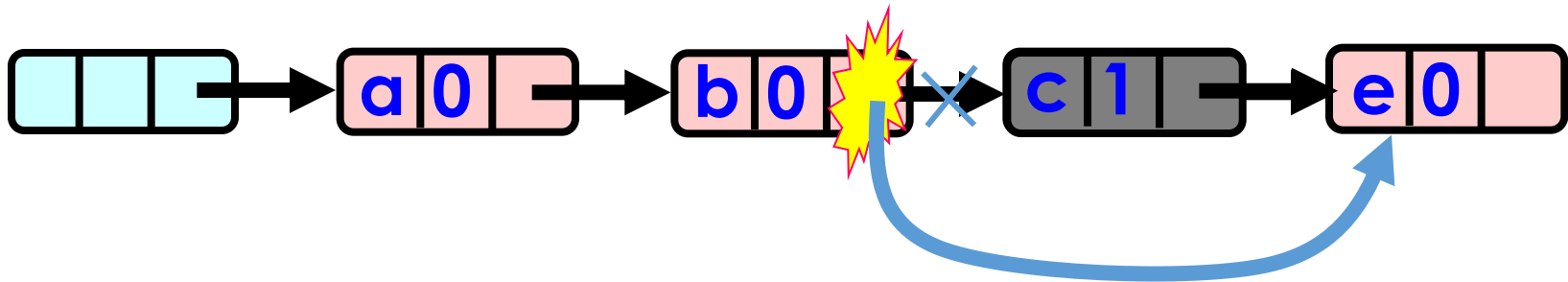
Lock-free Lists

- Next logical step (*after the lazy list*) is...
- Eliminate locking entirely
 - *contains()* wait-free
 - *add()* lock-free
 - *remove()* lock-free
- Use only *compareAndSet()*
- What could go wrong?

Remove Using CAS

- `remove(c)`

Logical Removal =
Set Mark Bit



Use CAS to verify
if pointer is correct

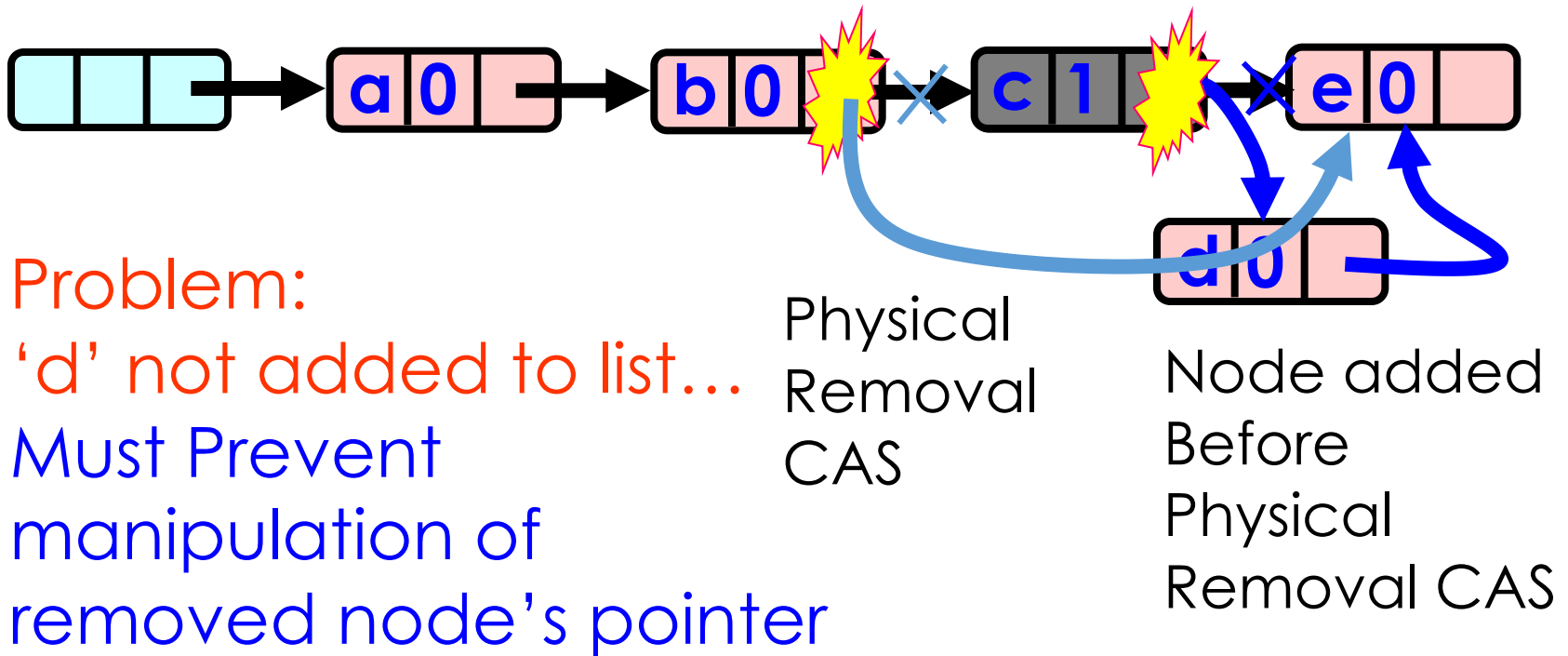
Physical
Removal
CAS pointer

Not enough! Why?

Problem...

- `remove(c) | add(d)`

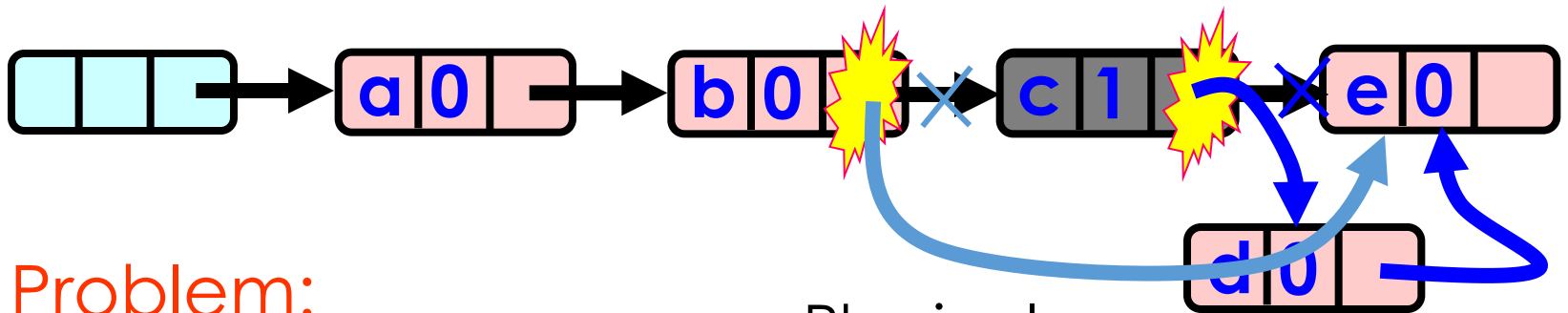
Logical Removal =
Set Mark Bit



Problem...

- `remove(c) | add(d)`

Logical Removal =
Set Mark Bit

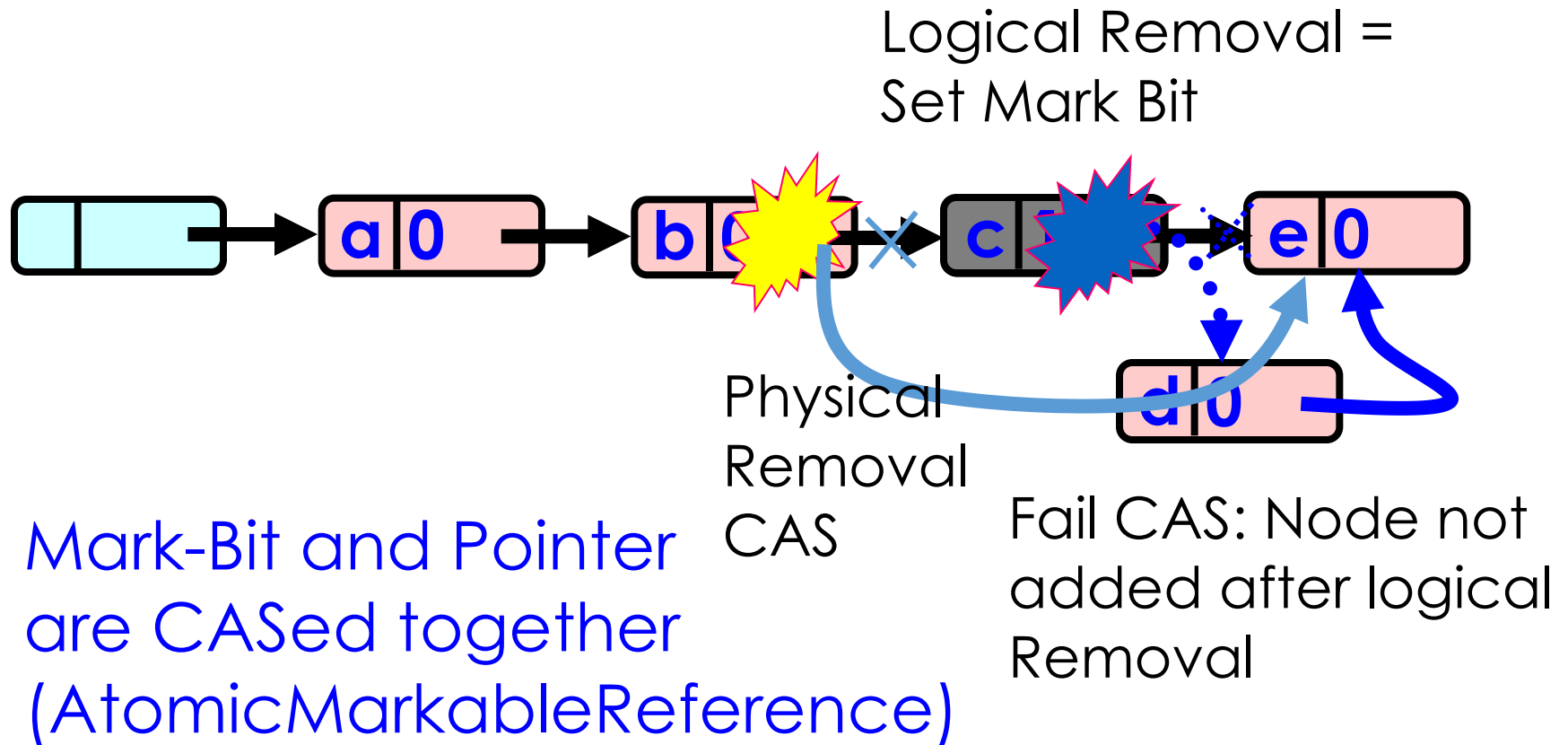


Problem:
'd' not added to list...
Must Prevent
manipulation of
removed node's pointer

Physical
Removal
CAS

Node added
Before
Physical
Removal CAS

The Solution: Combine Bit and Pointer



Solution

- Use *AtomicMarkableReference*
- Atomically
 - Swing reference and
 - Update flag
- Remove in two steps
 - Set mark bit in next field
 - Redirect predecessor's pointer with a CAS

Marking a Node

- AtomicMarkableReference class
 - java.util.concurrent.atomic package



Extracting Reference & Mark

```
public Object get(boolean[] marked);
```


Extracting Reference & Mark

```
public object get(boolean[] marked);
```

**Returns
reference**

**Returns mark at
array index 0!**

Extracting Reference Only

```
public boolean isMarked();
```

Value of
mark

Changing State

```
public boolean compareAndSet(  
    Object expectedRef,  
    Object updateRef,  
    boolean expectedMark,  
    boolean updateMark);
```

Changing State

□

If this is the current
reference ...

```
public boolean compareAndSet(  
    Object expectedRef,  
    Object updateRef,  
    boolean expectedMark,  
    boolean updateMark);
```

And this is the
current mark ...

Changing State

□

...then change to this
new reference ...

```
public boolean compareAndSet(  
    Object expectedRef,  
    Object updateRef,  
    boolean expectedMark,  
    boolean updateMark);
```

... and this new
mark

Changing State

```
public boolean attemptMark(  
    Object expectedRef,  
    boolean updateMark);
```

Changing State

```
public boolean attemptMark(  
    Object expectedRef,  
    boolean updateMark);
```

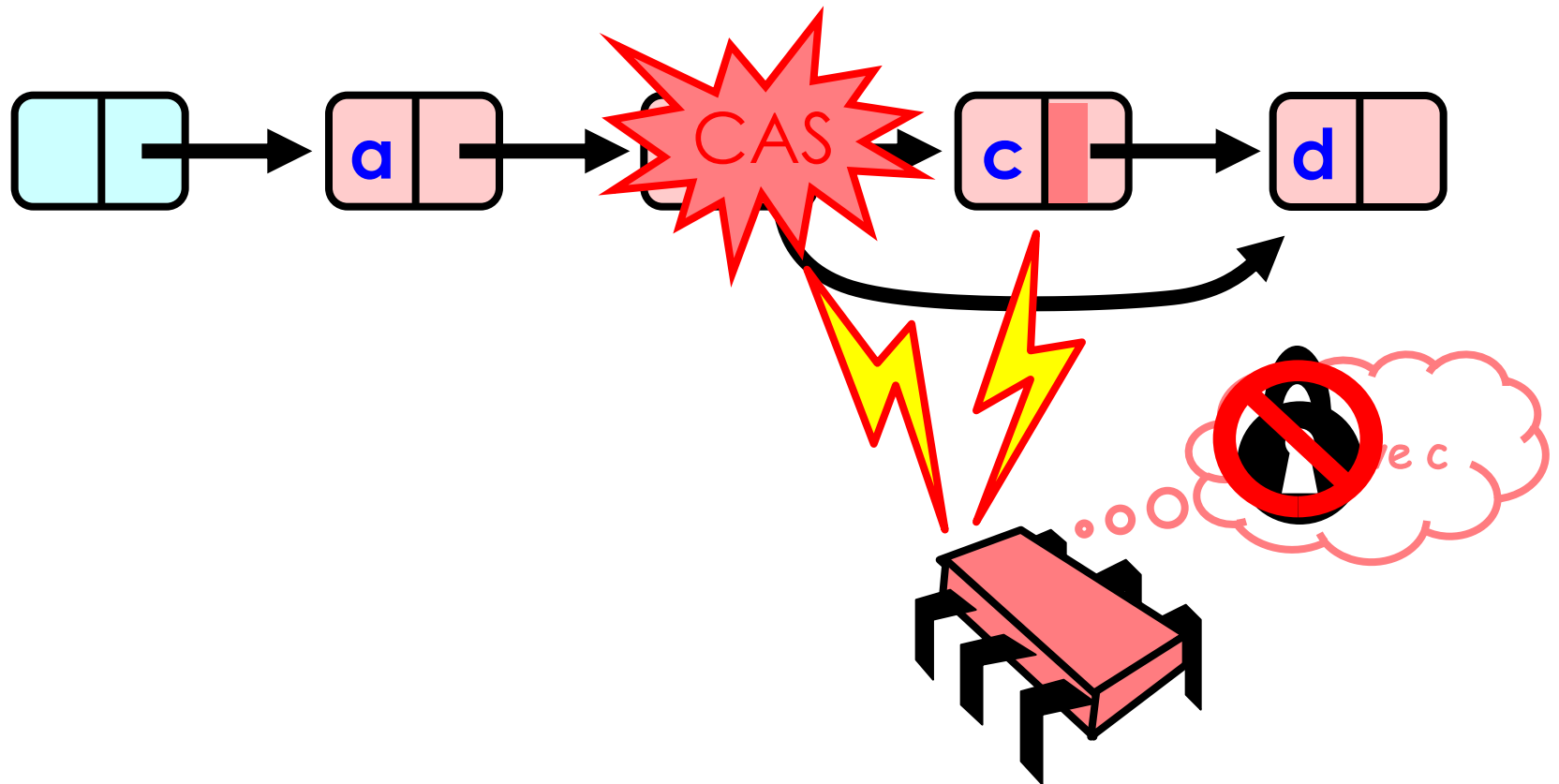
**If this is the current
reference ...**

Changing State

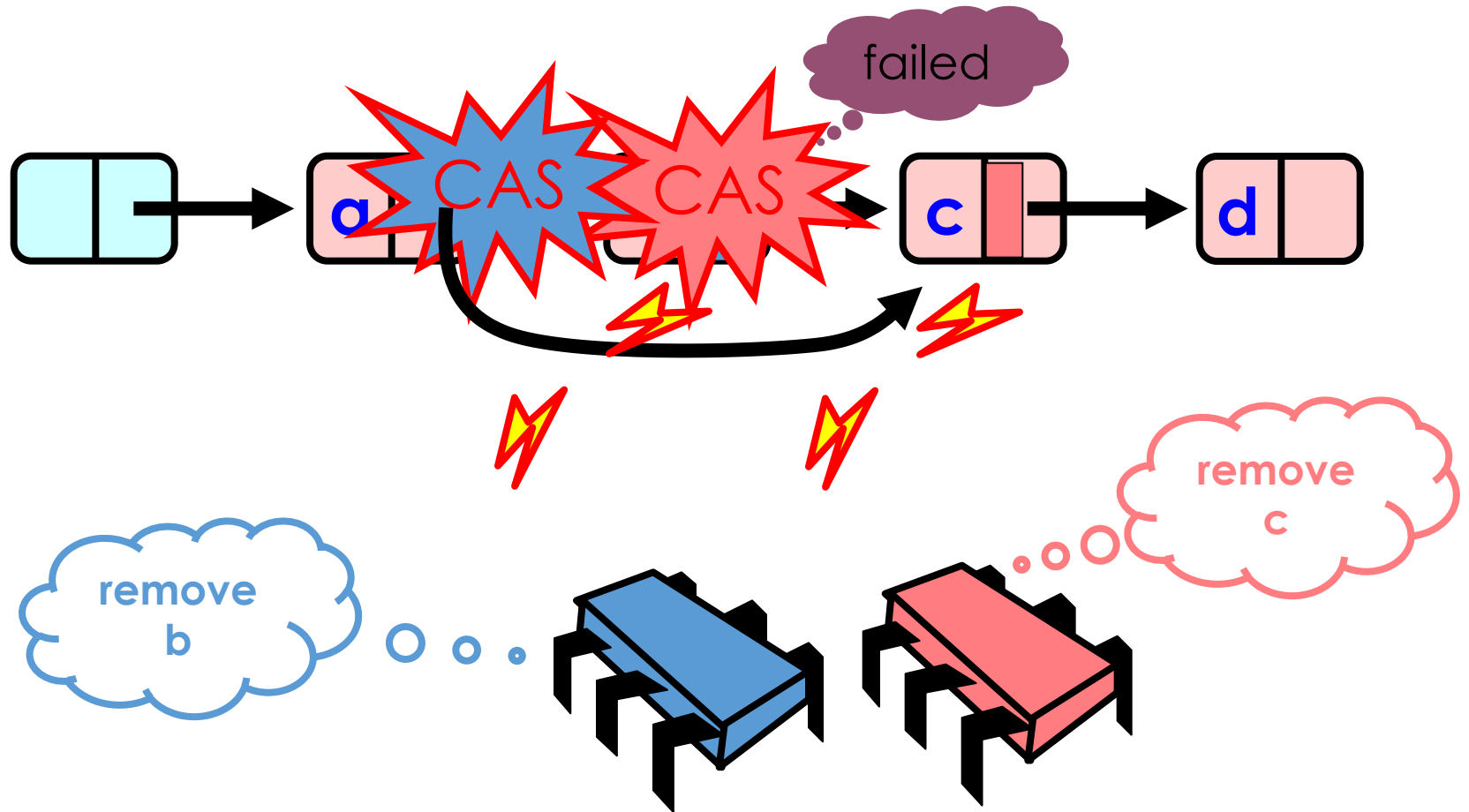
```
public boolean attemptMark(  
    Object expectedRef,  
    boolean updateMark);
```

**.. then change to
this new mark.**

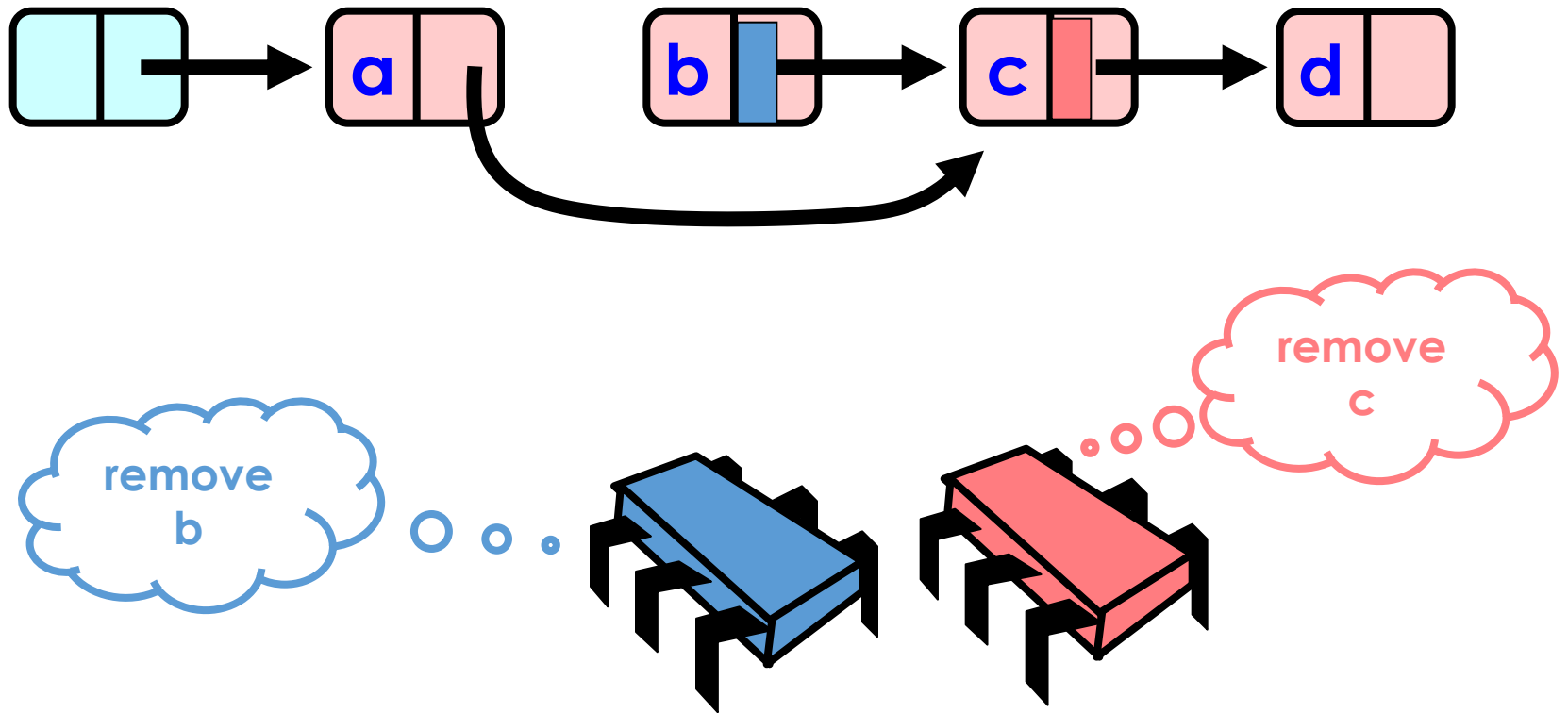
Removing a Node



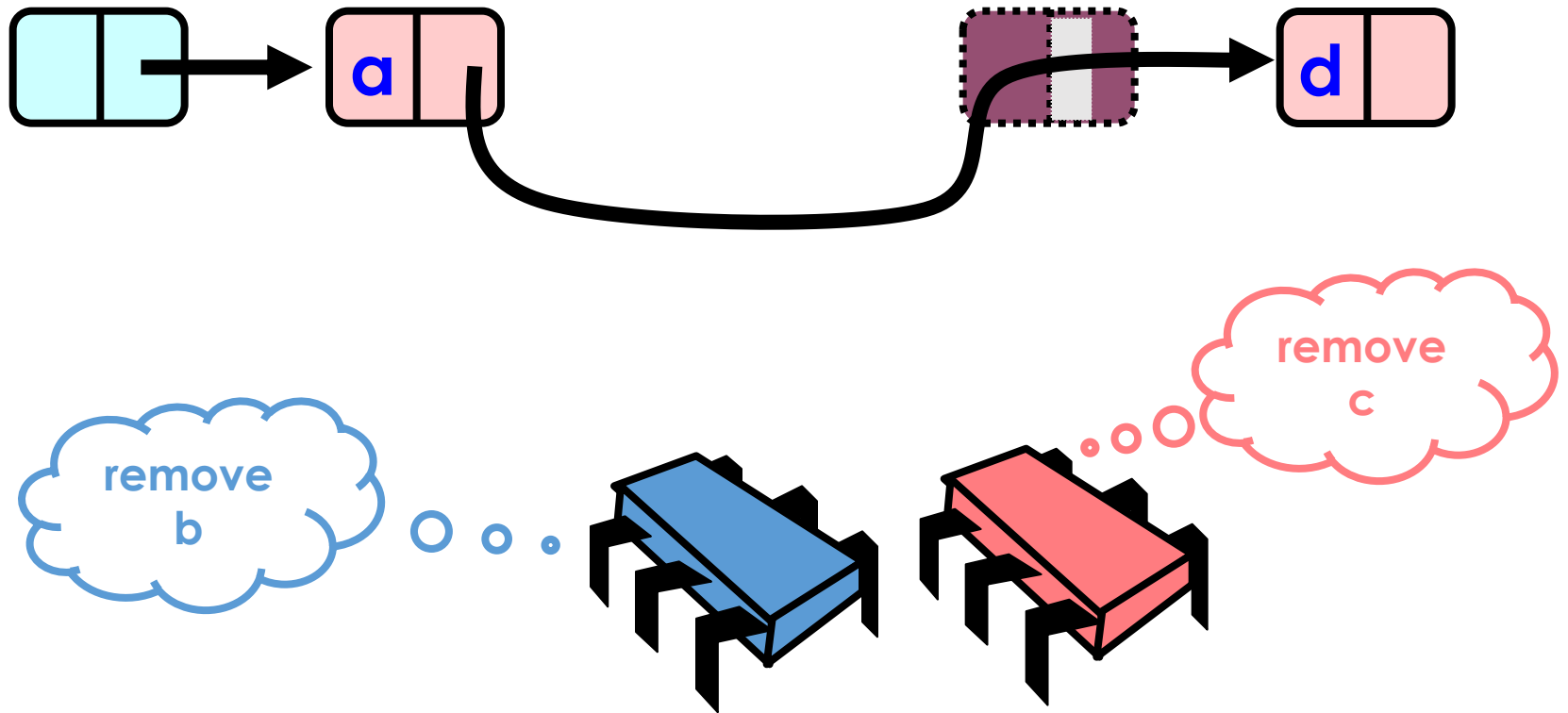
Removing a Node



Removing a Node



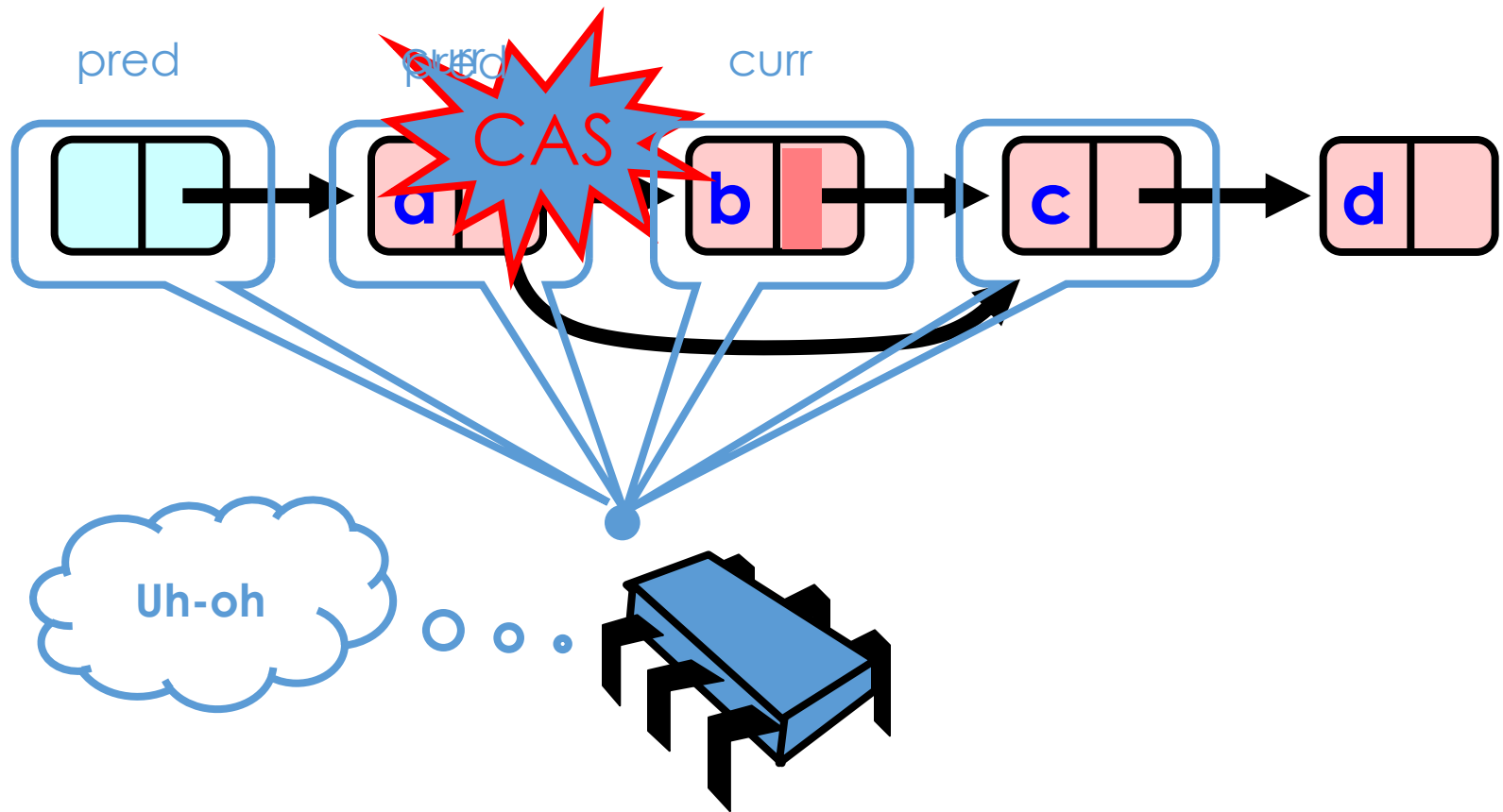
Removing a Node



Traversing the List

- Q: what do you do when you find a “logically” deleted node in your path?
- A: finish the job.
 - CAS the predecessor’s next field
 - Proceed (repeat as needed)

Lock-Free Traversal (only Add and Remove)



The END
